Solving equations in sequences: cans and cannots

Gleb Pogudin (LIX, CNRS, École Polytechnique, IPP) joint with A. Ovchinnikov, T. Scanlon, and M. Wibmer



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In this talk:

- algorithm for checking consistency of a system of equations (and elimination),
- undecidability results for almost anything beyond,
- and speculation.

Part 1: Prologue Main characters and first obstacles

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- has a solution $f = \{\dots, 0, 1, 0, 1, \dots\}$ is σ is a shift on $\mathbb{C}^{\mathbb{Z}}$.

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- in sequences: g contains at least one zero;
- in germs: g contains infinitely many zeros.

Not an easy solution space!

Theorem (Hrushovski, Point, 2007)

Problem:

- given a system of difference equations and inequations
- check if it has a solution.

Is undecidable both in $\mathbb{C}^{\mathbb{Z}}$ and \mathcal{G} .

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Glimpse under the hood

Encoding diophantine equations:

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implies that f consists of integers. Then

$$h = \sigma(h) \& (h - f) \cdot e = 0 \& e \neq 0$$

implies that h is a constant integer sequence.

Part 2: Cans

Consistency and elimination

joint with A. Ovchinnikov and T. Scanlon https://arxiv.org/abs/1712.01412 Given a system of difference equations over C
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Question: solution where?





Theorem (Ovchinnikov, Pogudin, Scanlon, 2020)

System of difference equations over a constant field k has \implies It has a solution in $\bar{k}^{\mathbb{Z}}$ a solution in some difference ring

Does there exist a sequence $\{a_n\}_{n\in\mathbb{Z}}$ such that:

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NO because

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Idea: no "finite" solution \implies no solution. Converse? Bound?

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Involves

- parts of the proof of bound (coming soon);
- nonstandard Frobenius as a model of ACFA.
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Possible approach

- 1. Apply shift $n \mapsto n+1$ (a prolongation) to the system
- 2. Check consistency of the polynomial system
- 3. If not succeed, go to Step 1.

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Half-solution: can detect inconsistency but not consistency.

Theorem (OPS, 2020)

If the system is inconsistent, this will be detected after at most

 $\mathbf{N}=B(d,D)$

prolongations, where

- **D** the degree of the system,
- **d** the dimension of the system.

$$B(d,D) = \begin{cases} D+1, \text{ if } d=0, \\ 2+D^2 + \frac{D(D-1)(D-2)}{6}, \text{ if } d=1, \\ B(d-1,D) + D^{B(d-1,D)}. \end{cases}$$

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This value is achieved on the elimination problem of x_i in

$$\begin{cases} x_{i+1} = x_i + 1, \\ x_i \cdot (x_i - 1) \dots \cdot (x_i - D + 1) = 0. \end{cases}$$

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Next idea: allow solutions to contain "high-dimensional points" For example: $\mathbb{A} = \mathbb{A} + 1$, so $\{a_n\} = \{\mathbb{A}, \mathbb{A}, \ldots\}$ is a periodic solution

Picture

We can bring every system to a form

$$(a_n + b_n)(a_n - b_n) = 0$$
 — nonlinear but no shifts

$$b_{n+1} = a_n$$
 — with a shift but linear

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- Sharpness in d = 0 and bound for d = 1 indicate that the worst case is a union of hyperplanes;
- Possible approach
 - consider the "worst" case, union of hypersurfaces, get lower bounds;
 - employ deformation argument to reduce general case to unions of hyperplanes.

Part 3: Cannots Implications, grids, ℝ joint with T. Scanlon and M. Wibmer https://arxiv.org/abs/1909.03239

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- Check if g = 0 holds for any solution of $f_1 = \ldots = f_{\ell} = 0$.

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 $\begin{array}{rcl} \mbox{Idea: system} \implies \mbox{piecie-wise polynomial map} \implies \\ \mbox{enumerating tuples of integers} \implies \mbox{diophantine equations} \end{array}$

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About the proof

In Hrushovski-Point: inequations used for " $\{a_n\}$ contains infinitely many zeroes". We do this in \mathbb{R} (with Lagrange four-square theorem!).

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Remark: similar result if the sequences are indexed by a free monoid.

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- Many interesting algebraically closed field: p-adics, $\mathbb{F}_p(t)$, etc
- Shink the class of sequences: the ones that may "come from discretization"?

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- New hope: algorithm for equations (+ bound, + universality, + elimination)
- Understanding the limits: undecidability for implications, reals, equations on grids
- Still many promising directions!
Thank you!

Looking for a PhD student: http: //www.lix.polytechnique.fr/Labo/Gleb.POGUDIN/phd-occam/